

Key Features Of A LAN

- High throughput
- Relatively low cost
- Limited to short distance
- Often rely on *shared media*

Scientific Justification For Local Area Networks

- *A computer is more likely to communicate with computers that are nearby than with computers that are distant.*

Known as the *locality principle*

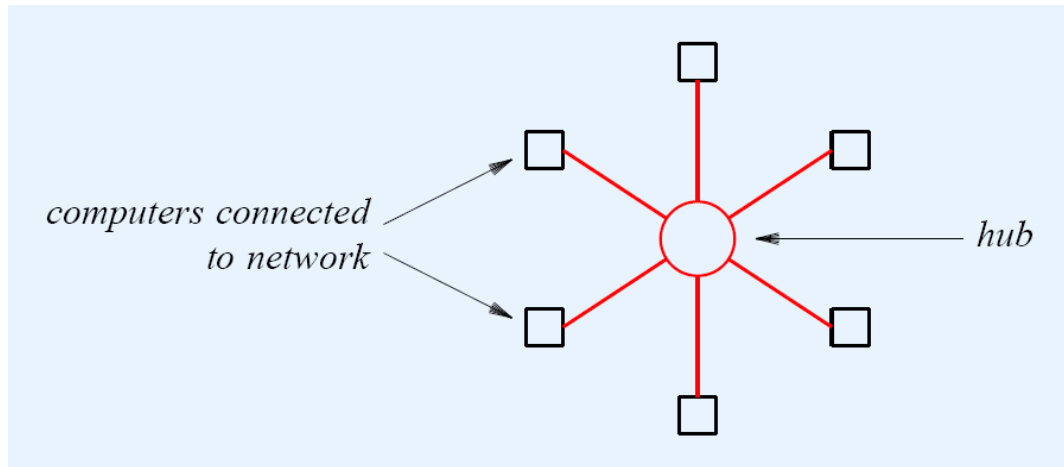
Topology

- Mathematical term
- Roughly interpreted as “geometry for curved surfaces”

Network Topology

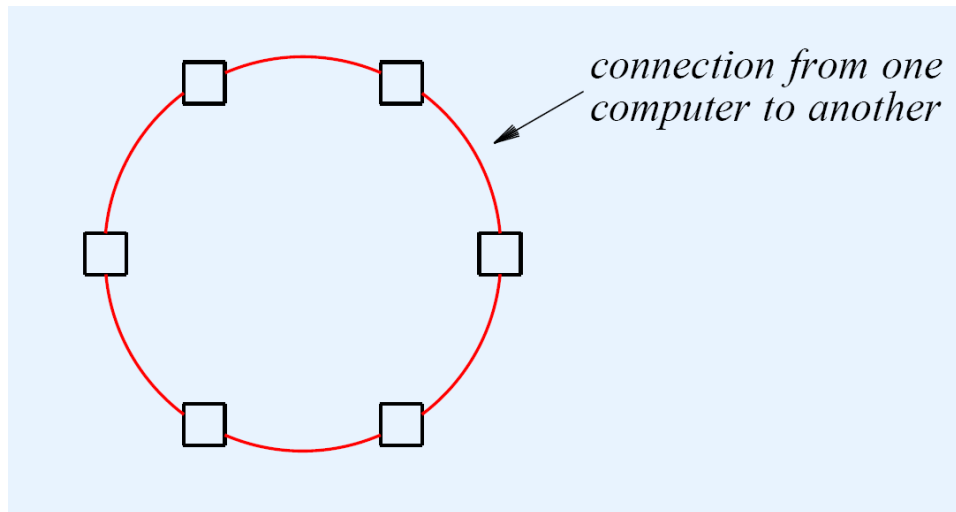
- Specifies general “shape” of a network
- Handful of broad categories
- Often applied to LAN
- Primarily refers to interconnections
- Hides details of actual devices

Star Topology



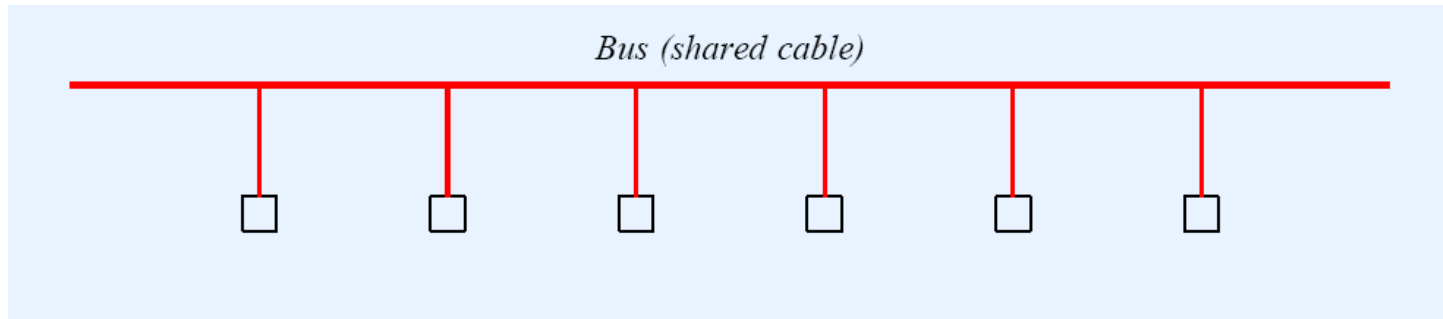
- Central component of network known as *hub*
- Each computer has separate connection to hub

Ring Topology



- No central facility
- Connections go directly from one computer to another

Bus Topology



- Shared medium forms main interconnect
- Each computer has a connection to the medium

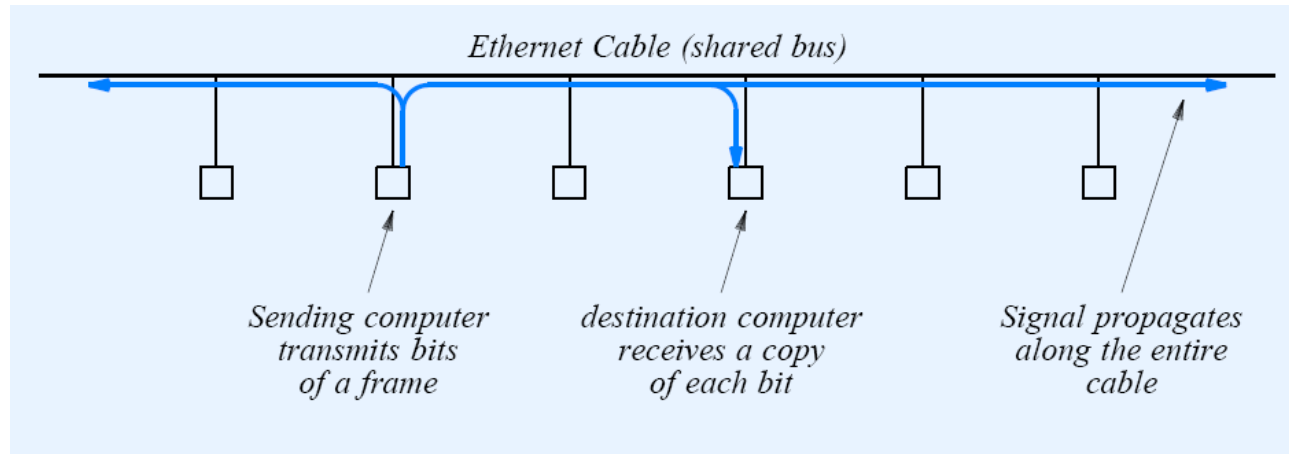
Example Bus Network: Ethernet

- Most popular LAN
- Widely used
- IEEE standard *802.3*
- Several generations
 - Same frame format
 - Different data rates
 - Different wiring schemes

Shared Medium In A LAN

- Shared medium used for all transmissions
- Only one station transmits at any time
- Stations “take turns” using medium
- Media Access Control (*MAC*) policy ensures fairness

Illustration Of Ethernet Transmission



- Only one station transmits at any time
- Signal propagates across entire cable
- All stations receive transmission
- CSMA/CD media access scheme

CSMA/CD Paradigm

- Multiple Access (*MA*)
 - Multiple computers attach to shared media
 - Each uses same access algorithm
- Carrier Sense (*CS*)
 - Wait until medium idle
 - Begin to transmit frame
- Simultaneous transmission possible

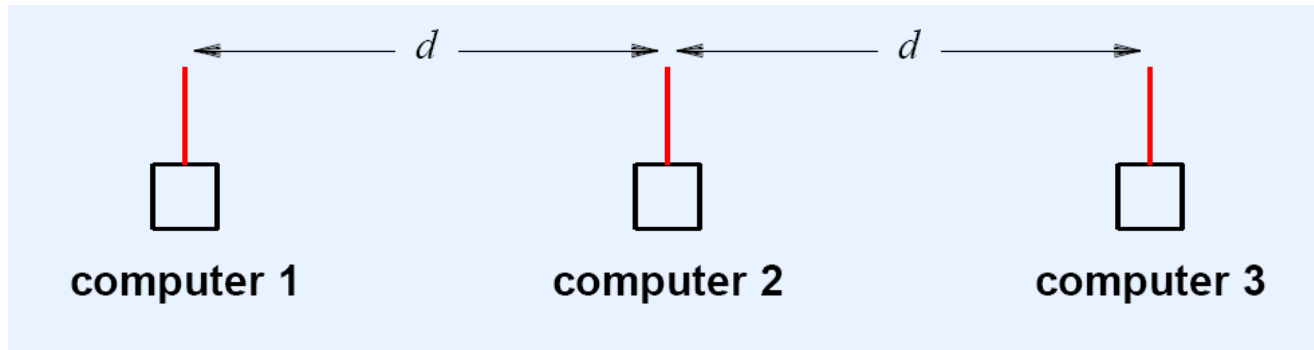
CSMA/CD Paradigm (continued)

- Two simultaneous transmissions
 - Interfere with one another
 - Called *collision*
- CSMA plus *Collision Detection (CD)*
 - Listen to medium during transmission
 - Detect whether another station's signal interferes
 - Back off from interference and try again

Backoff After Collision

- When collision occurs
 - Wait random time $t1$, $0 \leq t1 \leq d$
 - Use CSMA and try again
- If second collision occurs
 - Wait random time $t2$, $0 \leq t2 \leq 2d$
- Double range for each successive collision
- Called *exponential backoff*

Media Access On A Wireless Net



- Limited range
 - Not all stations receive all transmissions
 - Cannot use CSMA/CD
- Example in diagram
 - Maximum transmission distance is d
 - Stations 1 and 3 do not receive each other's transmissions

CSMA/CA

- Used on wireless LANs
- Both sides send small message followed by data transmission
 - “ X is about to send to Y ”
 - “ Y is about to receive from X ”
 - Data frame sent from X to Y
- Purpose: inform all stations in range of X or Y before transmission
- Known as *Collision Avoidance (CA)*

Wi-Fi Wireless LAN Technology

- Popular
- Uses CSMA/CA for media access
- Standards set by IEEE
 - 802.11b (11 Mbps, shared channel)
 - 802.11a (54 Mbps, shared channel)
- Named *Wi-Fi* by consortium of vendors (to enhance popular appeal)

Identifying A Destination

- All stations on shared-media LAN receive all transmissions
- To allow sender to specify destination
 - Each station assigned unique number
 - Known as station's *address*
 - Each frame contains address of intended recipient

Ethernet Addressing

- Standardized by IEEE
- Each station assigned unique 48-bit address
- Address assigned when network interface card (*NIC*) manufactured

Ethernet Address Recognition

- Each frame contains destination address
- All stations receive a transmission
- Station discards any frame addressed to another station
- Important: interface hardware, not software, checks address

Possible Destinations

- Packet can be sent to:
 - Single destination (*unicast*)
 - All stations on network (*broadcast*)
 - Subset of stations (*multicast*)
- Address used to distinguish

Advantages Of Address Alternatives

- Unicast
 - Efficient for interaction between two computers
- Broadcast
 - Efficient for transmitting to all computers
- Multicast
 - Efficient for transmitting to a subset of computers

Broadcast On Ethernet

- All 1s address specifies broadcast
- Sender
 - Places broadcast address in frame
 - Transmits one copy on shared network
 - All stations receive copy
- Receiver always accepts frame that contains
 - Station's unicast address
 - The broadcast address

Multicast On Ethernet

- Half of addresses reserved for multicast
- Network interface card
 - Always accepts unicast and broadcast
 - Can accept zero or more multicast addresses
- Software
 - Determines multicast address to accept
 - Informs network interface card

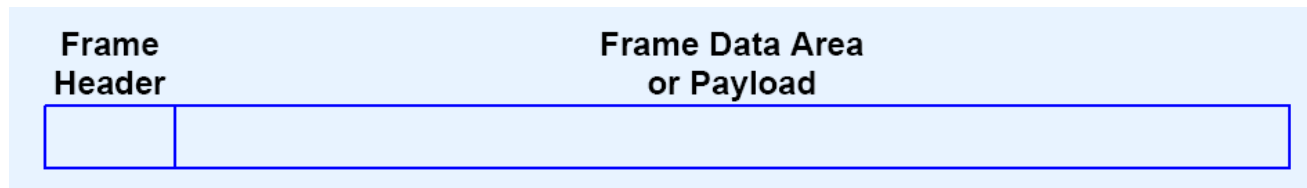
Promiscuous Mode

- Designed for testing/debugging
- Allows interface to accept *all* packets
- Available on most interface hardware

Identifying Frame Contents

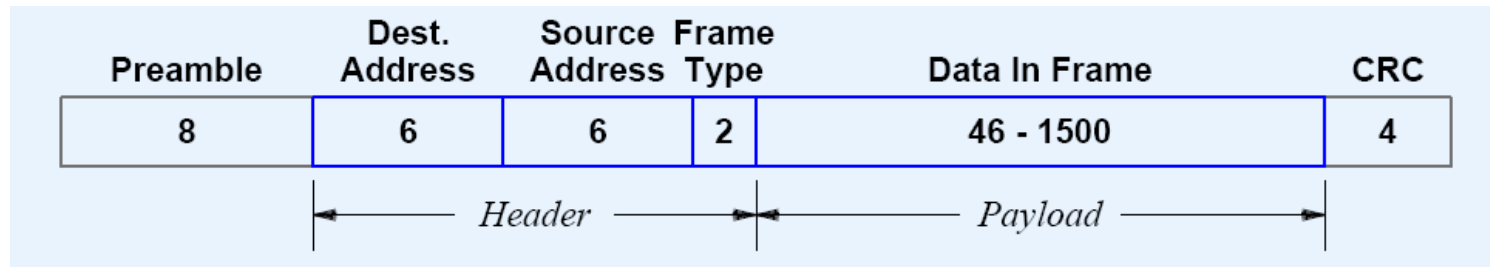
- Integer *type field* tells recipient the type of data being carried
- Two possibilities
 - *Self-identifying* or *explicit type* (hardware records type)
 - *Implicit type* (application sending data must handle type)

Conceptual Frame Format



- *Header*
 - Contains address and type information
 - Layout fixed
- *Payload*
 - Contains data being sent

Illustration Of Ethernet Frame



- Sender places
 - Sender's address in *source*
 - Recipient's address in *destination*
 - Type of data in *frame type*
 - Cyclic redundancy check in *CRC*

Example Ethernet Types

Value

0000–05DC

0800

0805

0900

0BAD

1000–100F

6004

6559

8005

8008

8014

8035

8038

805C

809B

80C4–80C5

80D5

80FF–8103

8137–8138

818D

FFFF

Meaning

Reserved for use with IEEE LLC/SNAP

Internet IP Version 4

CCITT X.25

Ungermann–Bass Corporation network debugger

Banyan Systems Corporation VINES

Berkeley UNIX Trailer encapsulation

Digital Equipment Corporation LAT

Frame Relay

Hewlett Packard Corporation network probe

AT&T Corporation

Silicon Graphics Corporation network games

Internet Reverse ARP

Digital Equipment Corporation LANBridge

Stanford University V Kernel

Apple Computer Corporation AppleTalk

Banyan Systems Corporation

IBM Corporation SNA

Wellfleet Communications

Novell Corporation IPX

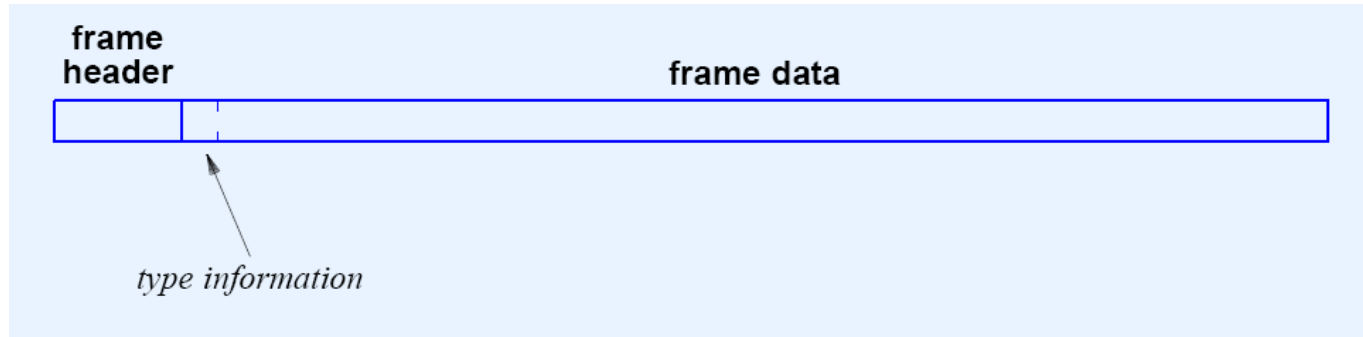
Motorola Corporation

Reserved

When Network Hardware Does Not Include Types

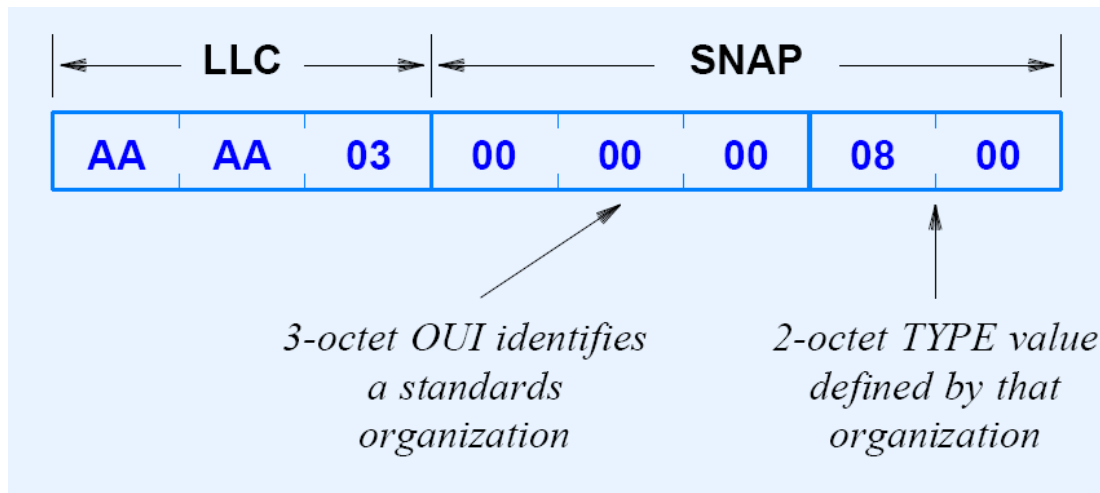
- Sending and receiving computers must agree
 - To only send one type of data
 - To put type information in first few octets of payload
- Most systems need type information

Illustration Of Type Information Added To Data



- In practice
 - Type information small compared to data carried
 - Format of type information standardized

A Standard For Type Information



- Defined by IEEE
- Used when hardware does not include type field
- Called *LLC/SNAP header*

Demultiplexing On Type

- Network interface hardware
 - Receives copy of each transmitted frame
 - Examines address and either discards or accepts
 - Passes accepted frame to system software
- Network device software
 - Examines frame type
 - Passes frame to correct software module

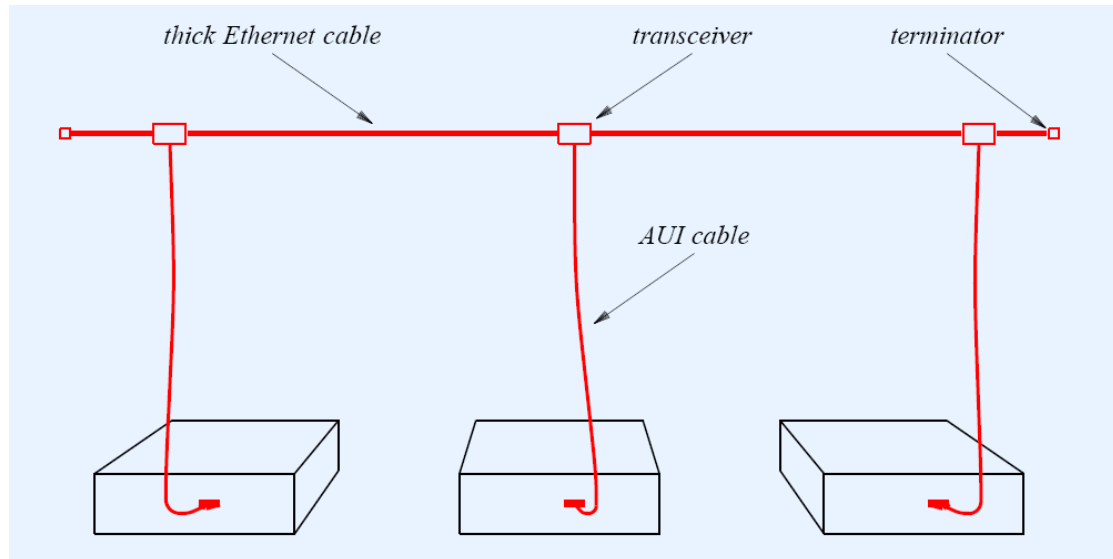
Network Analyzer

- Device used for testing and maintenance
- Listens in promiscuous mode
- Produces
 - Summaries (e.g., % of broadcast frames)
 - Specific items (e.g., frames from a given address)

Ethernet Wiring

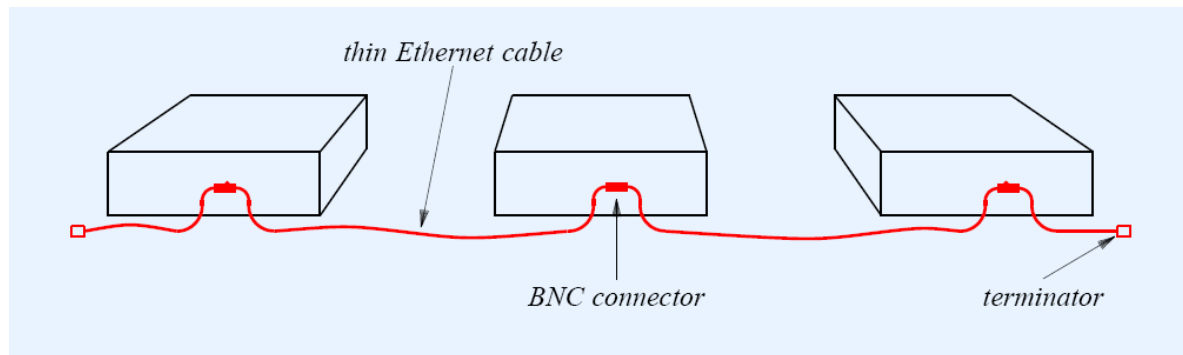
- Three schemes
 - Correspond to three generations
 - All use same frame format

Original Ethernet Wiring



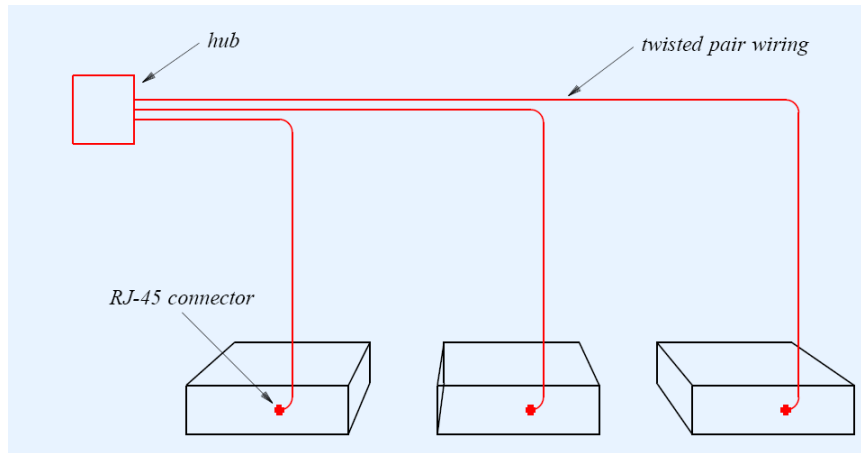
- Used heavy coaxial cable
- Formal name *10Base5*
- Called *thicknet*

Second Generation Ethernet Wiring



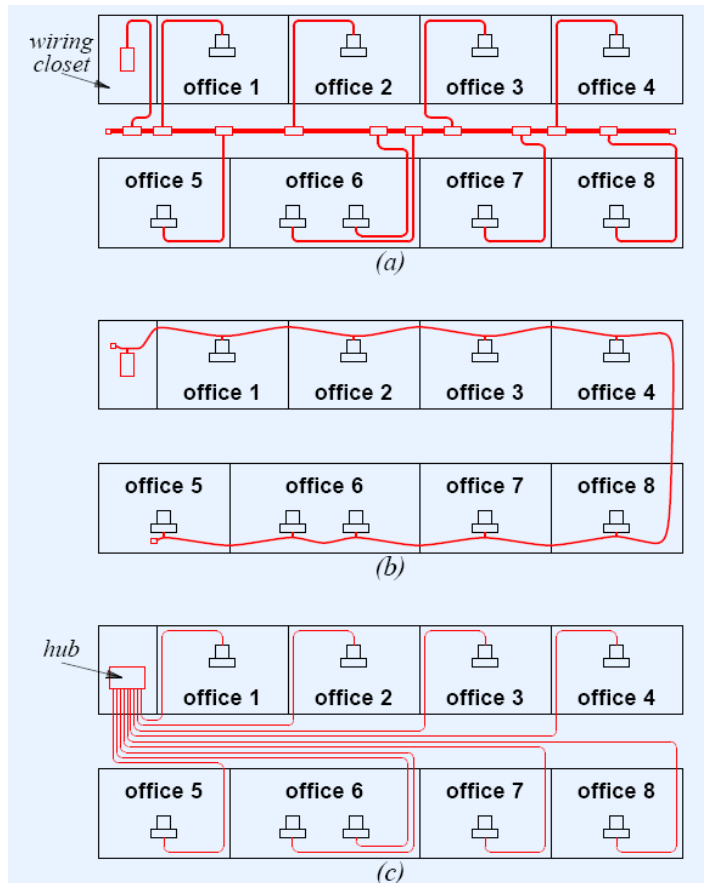
- Used thinner coaxial cable
- Formal name *10Base2*
- Called *thinnet*

Modern Ethernet Wiring



- Uses a *hub*
- Formal name *10Base-T*
- Called *twisted pair Ethernet*

Ethernet Wiring In An Office



A Note About Ethernet Topology

- Apparently
 - Original Ethernet used bus topology
 - Modern Ethernet uses star topology
- In fact, modern Ethernet is
 - Physical star
 - Logical bus
 - Called *star-shaped bus*

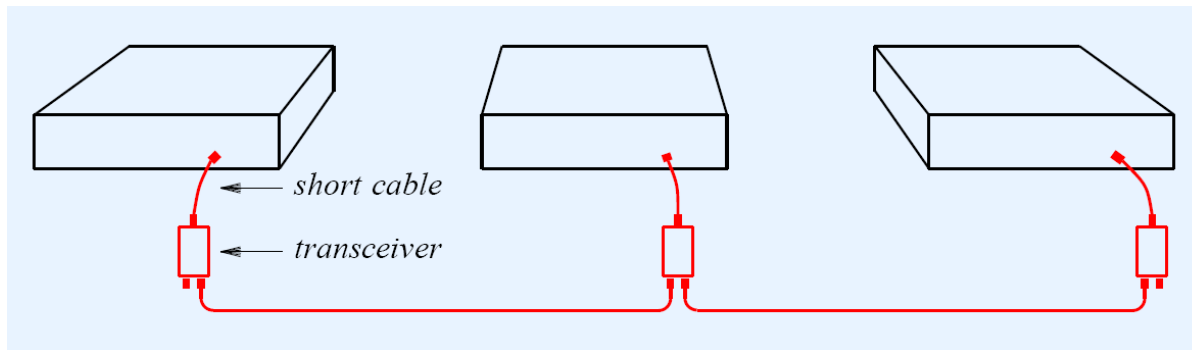
Higher Speed Ethernets

- *Fast Ethernet*
 - Operates at 100 Mbps
 - Formally *100Base-T*
 - Two wiring standards
 - 10/100 Ethernet devices available
- *Gigabit Ethernet*
 - Operates at 1000 Mbps (1 Gbps)
 - Slightly more expensive

Another LAN Using Bus Topology

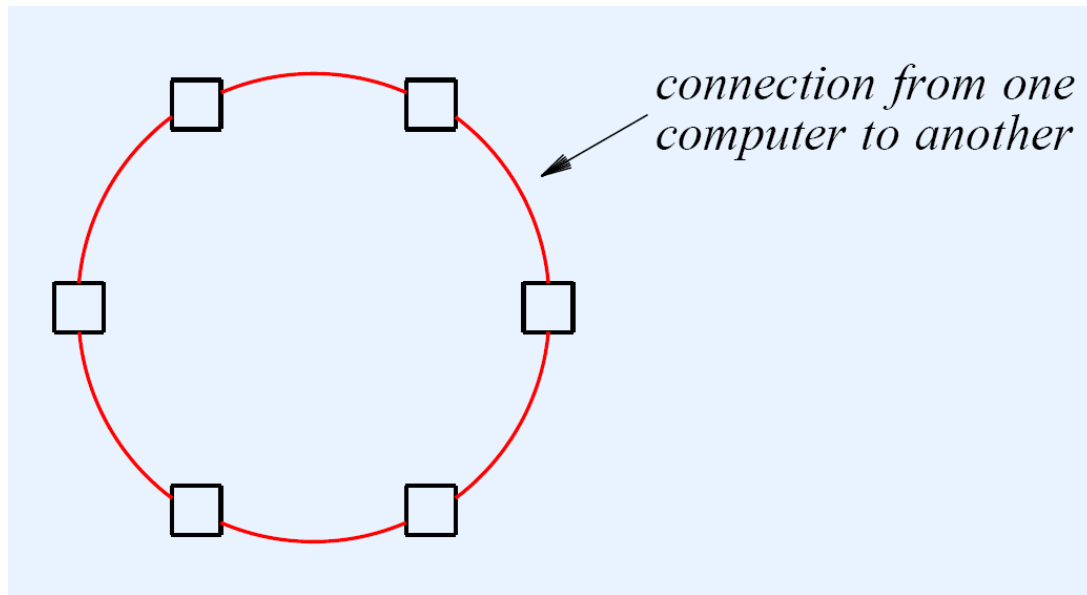
- LocalTalk
 - Developed by Apple Corp.
 - Simple to use
 - Slow by current standards

Illustration Of LocalTalk



- Transceiver required per station
- Transceiver terminates cable

Ring Topology



- Once a popular topology for LANs
- Bits flow in single direction

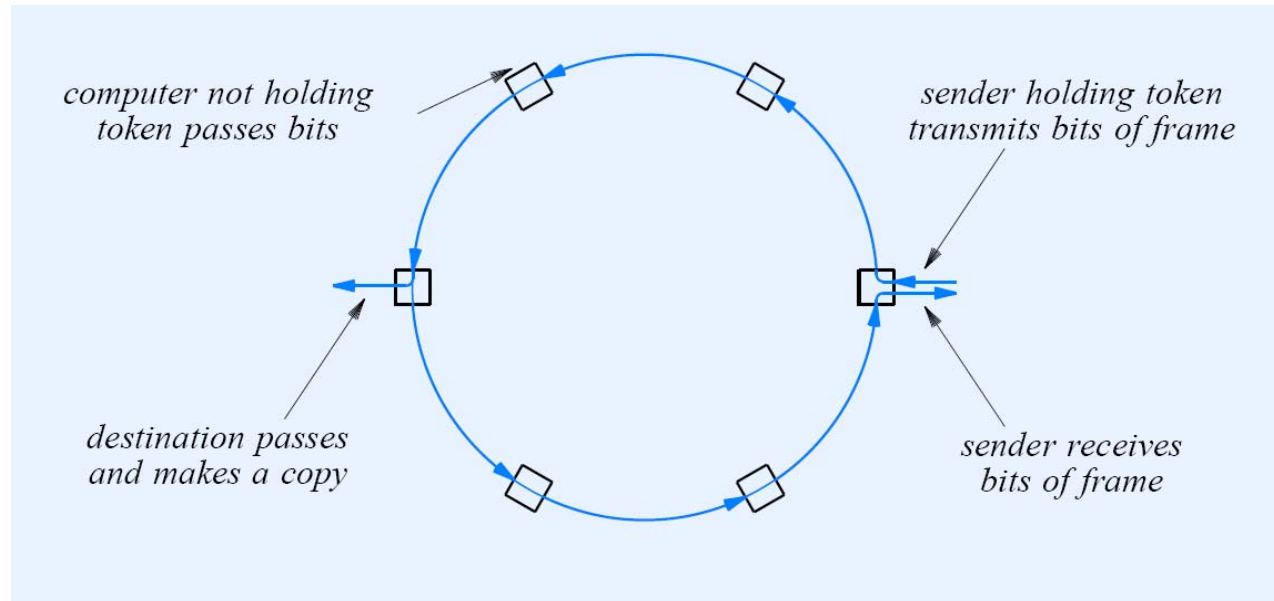
Token Passing

- Designed for ring topology
- Guarantees fair access
- *Token*
 - Special (reserved) message
 - Small (a few bits)

Token Passing Paradigm

- Station
 - Waits for token to arrive
 - Transmits one packet around ring
 - Transmits token around ring
- When no station has data to send
 - Token circulates continuously

Token Passing Ring Transmission



- Station waits for token before sending
- Signal travels around entire ring
- Sender receives its own transmission

Strengths Of Token Ring Approach

- Easy detection of
 - Broken ring
 - Hardware failures
 - Interference

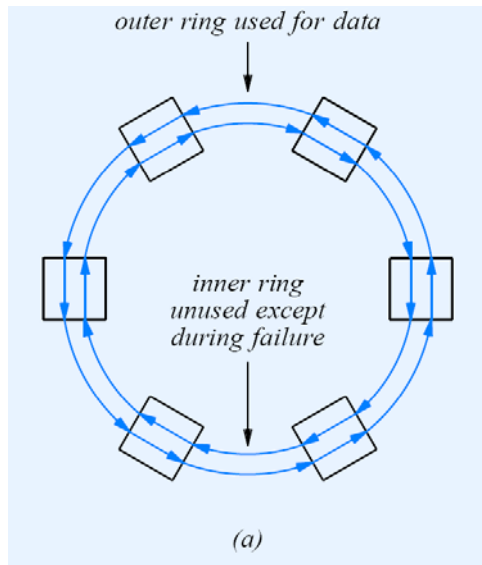
Weaknesses Of Token Ring Approach

- Broken wire disables entire ring
- Point-to-point wiring
 - Awkward in office environment
 - Difficult to add/move stations

Failure Recovery In Ring Networks

- Automatic failure recovery
- Introduced by FDDI
- Uses two rings
- Terminology
 - *Dual-attached*
 - *Counter rotating*
 - *Self healing*

Illustration Of Failure Recovery



- Normal operation uses one of two rings