

Need For Protocols (continued)

- Need mechanisms to distinguish among
 - Multiple computers on a network
 - Multiple applications on a computer
 - Multiple copies of a single application on a computer

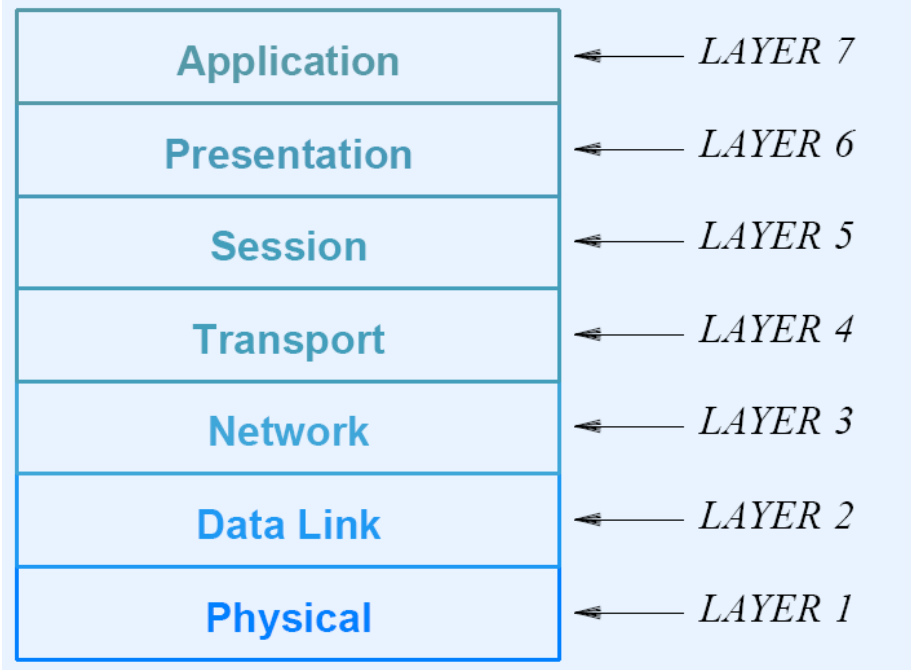
Set Of Protocols

- Work together
- Each protocol solves part of communication problem
- Known as
 - Protocol suite
 - Protocol family
- Designed in layers

Plan For Protocol Design

- Intended for protocol designers
- Divides protocols into layers
- Each layer devoted to one subproblem
- Example: ISO 7-layer reference model

Illustration Of The 7-Layer Model



- Defined early (by ISO (International Organization for Standardization))
- Now somewhat dated
- Does not include Internet layer!

ISO Layers

- Layer 1: Physical
 - Underlying hardware (e.g. RS232)
- Layer 2: Data Link (media access)
 - Hardware frame definitions (testing) + (transmission)
(e.g. frame format, byte stuffing, checksum)
- Layer 3: Network
 - Packet forwarding (e.g. IP addressing)
- Layer 4: Transport
 - Reliability (e.g. UDP/TCP)

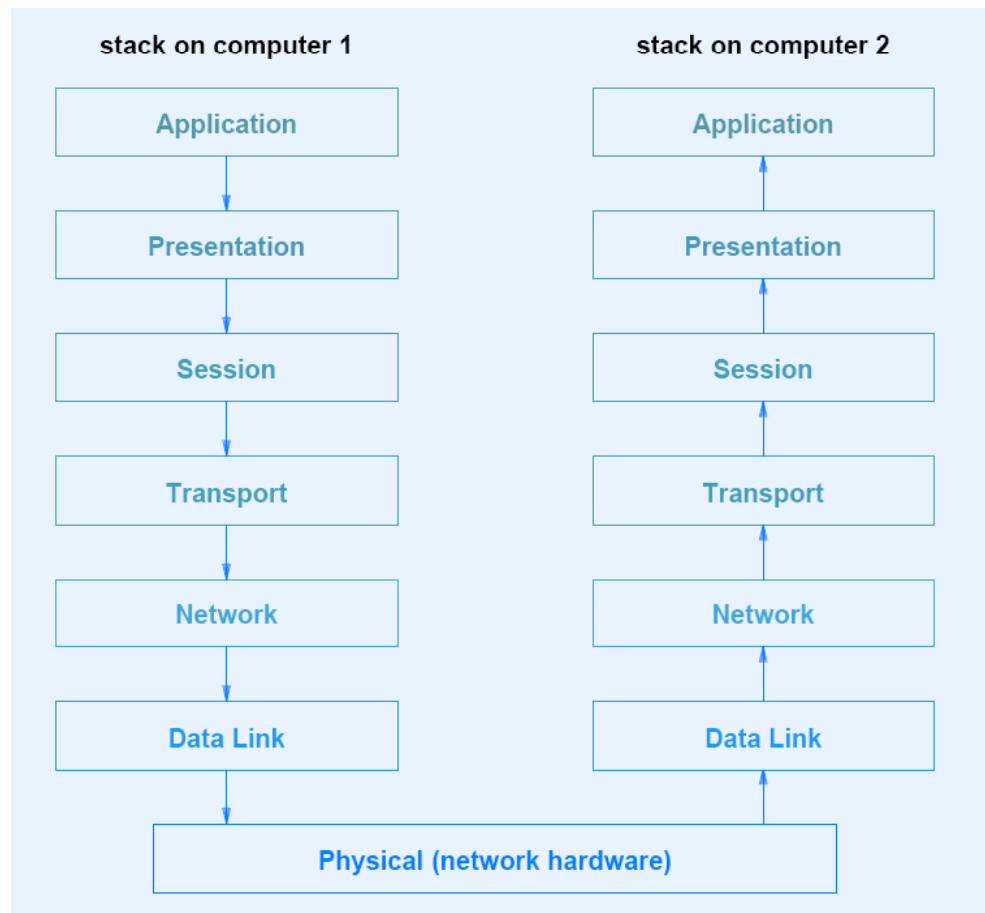
ISO Layers (continued)

- Layer 5: Session
 - How to establish a communication session (e.g. Login and passwords, authentication)
- Layer 6: Presentation
 - Data representation (differing brands uses differing representations of integers and characters)
- Layer 7: Application
 - Individual application program

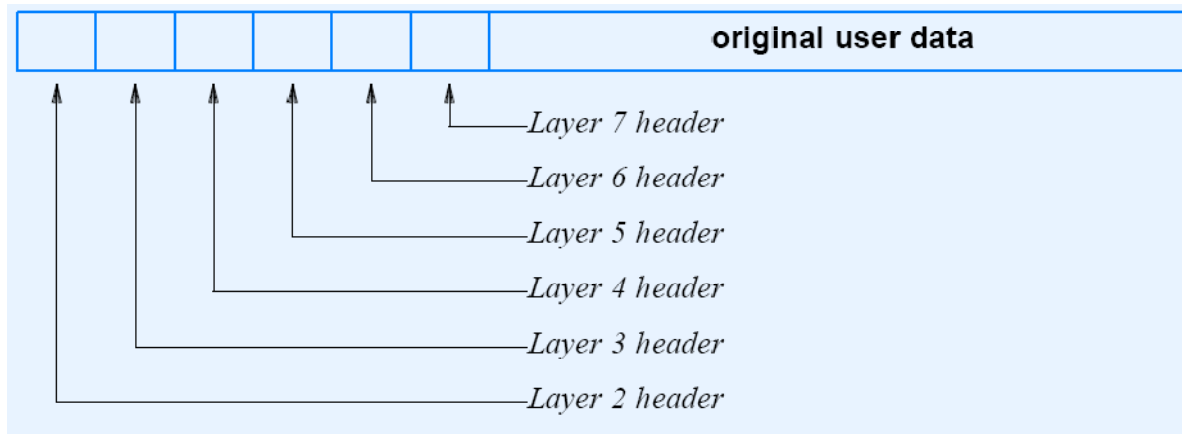
Layers And Protocol Software

- Protocol software follows layering model
 - One software module per layer
 - Modules cooperate
 - Incoming or outgoing data passes from one module to another
- Entire set of modules known as stack

Illustration Of Stacks



Layers And Packet Headers



- Lower level protocol header first

- No Header for layer 1

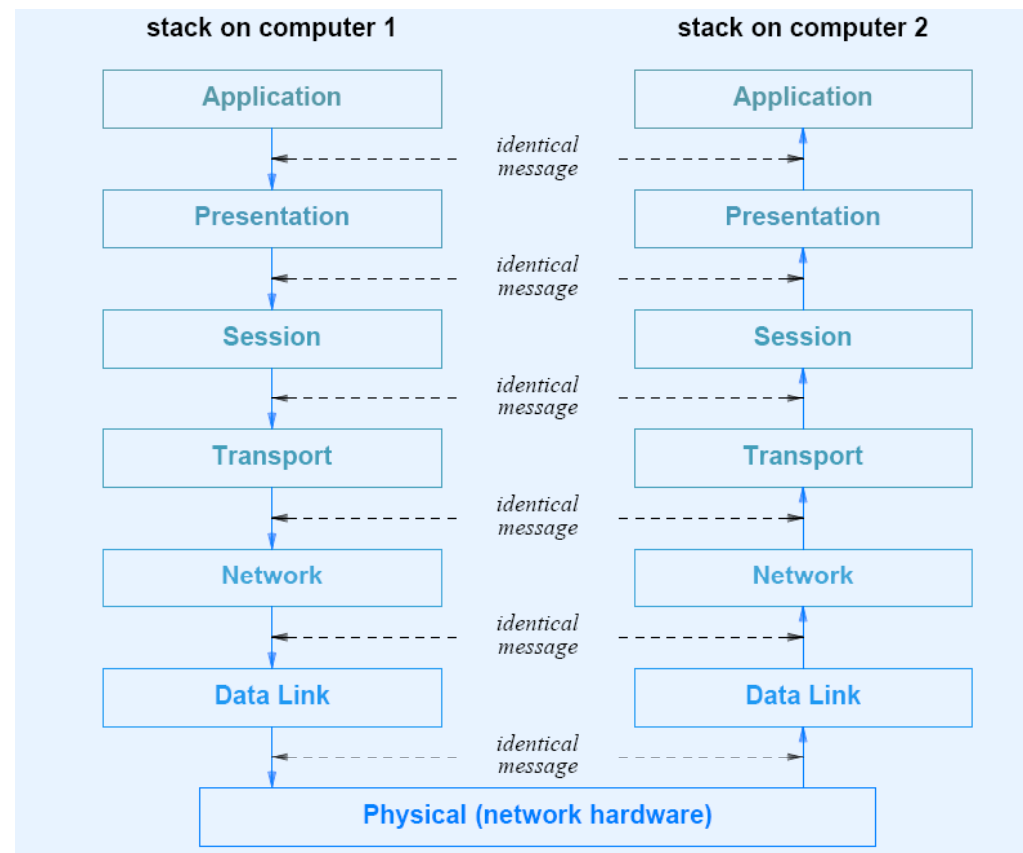
- Each layer
 - Prepends header to outgoing packet (some protocols appends trailer, too)
 - Removes header from incoming packet

Scientific Layering Principle

Software implementing layer N at the destination receives exactly the message sent by software implementing layer N at the source.

Change in a layer is invisible to other layers,
Can be designed, implemented, tested
independently w.r.t other layers

Illustration Of Layering Principle



Protocol Techniques

(Various techniques for all the layers)

- For bit corruption
 - Parity
 - Checksum
 - CRC
- For out-of-order delivery
(usually connectionless network)
 - (i+1)th by shorter path compared to ith packet
 - Sequence numbers
- Duplication
(e.g. sender transmits twice during collision in CSMA/CD)
 - Sequence numbers

Protocol Techniques (continued)

- For lost packets
 - Positive acknowledge and retransmission
(Set a timer to measure ACK response)
(Should handle duplicate packet by retransmitting lost packets
– but too early decision causes the duplication)
- For replay (excessive queuing delay)
 - Unique message ID (for each session)
 - (Previous session terminated if new session starts for replay.)
(e.g.: time the session is established)
- For data overrun
 - Flow control
(e.g.: stop-and-go (single packet basis) → inefficient)
sliding window (window size basis))

Flow Control

- Needed because
 - Sending computer faster than receiving computer
 - Sending application faster than receiving application
- Related to buffering
- Two forms
 - Stop-and-go
 - Sliding window

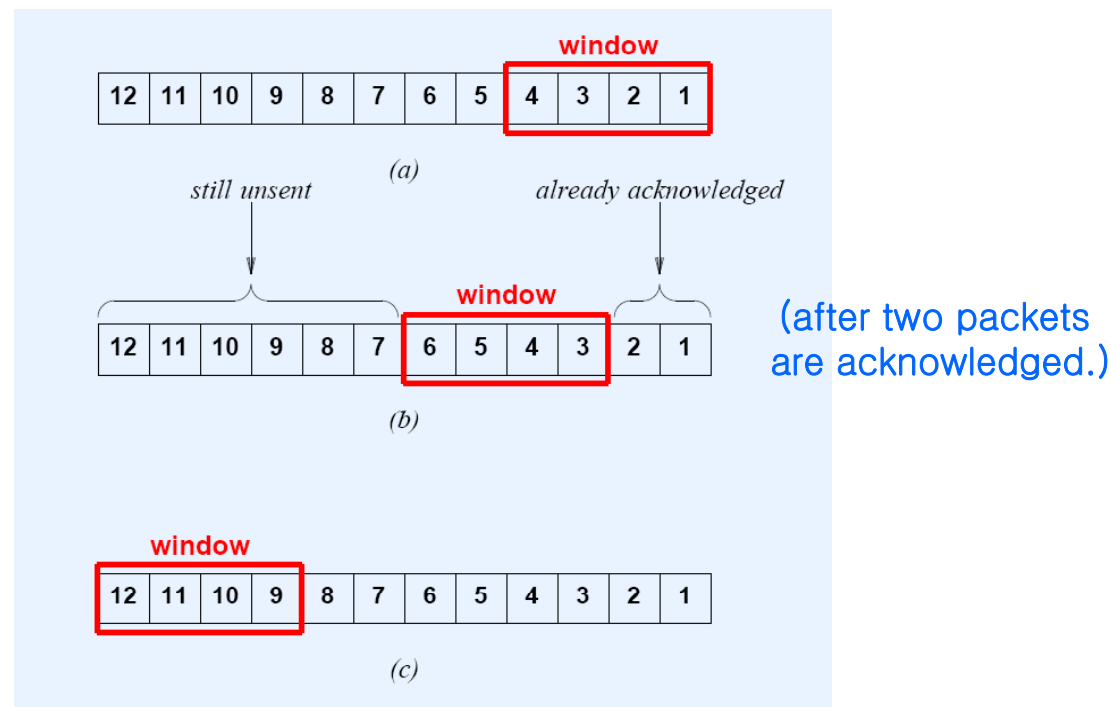
Stop-And-Go Flow Control

- Sending side
 - Transmits one packet
 - Waits for signal from receiver
- Receiving side
 - Receives and consumes packet
 - Transmits signal to sender
- Inefficient

Sliding Window Flow Control

- Receiving side
 - Establishes multiple buffers and informs sender
- Sending side
 - Transmits packets for all available buffers (@receiver)
 - Only waits if no signal arrives before transmission completes
- Receiving side
 - Sends signals as packets arrive

Illustration Of Sliding Window On Sending Side

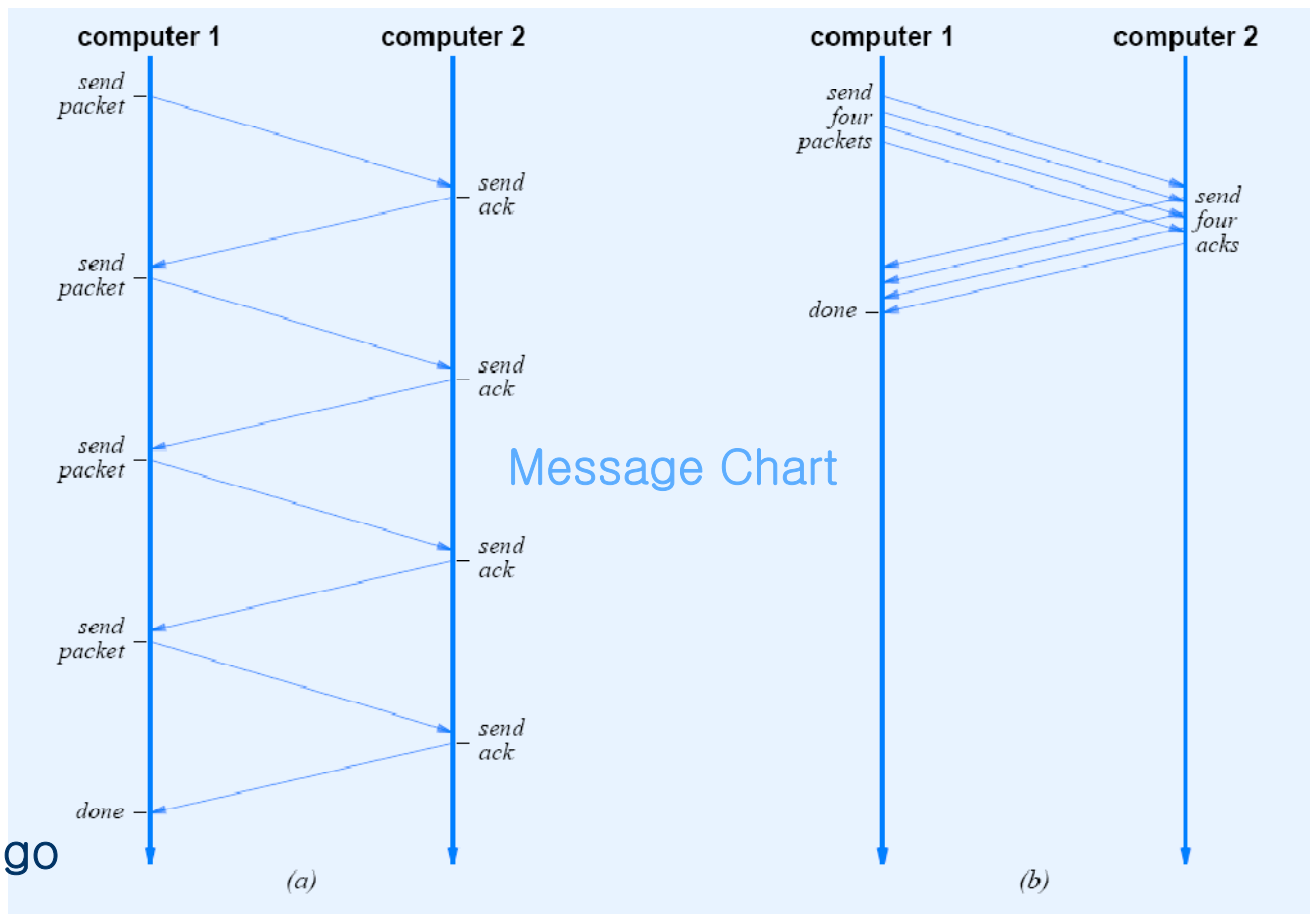


- Window tells how many packets can be sent
- Window moves as acknowledgements arrive

Performance

- Stop-and-go
 - Slow
 - Useful only in special cases
- Sliding window
 - Fast
 - Needed in high-speed network

Comparison Of Flow Control



Stop and go

Sliding Window

Why Sliding Window?

- Simultaneously
 - Increase throughput
 - Control flow
- Speedup

$$T_w = \min (B, T_g \times W)$$

where

- * T_w is sliding window throughput
- * B is underlying hardware bandwidth
- * T_g is stop-and-go throughput
- W is window size.

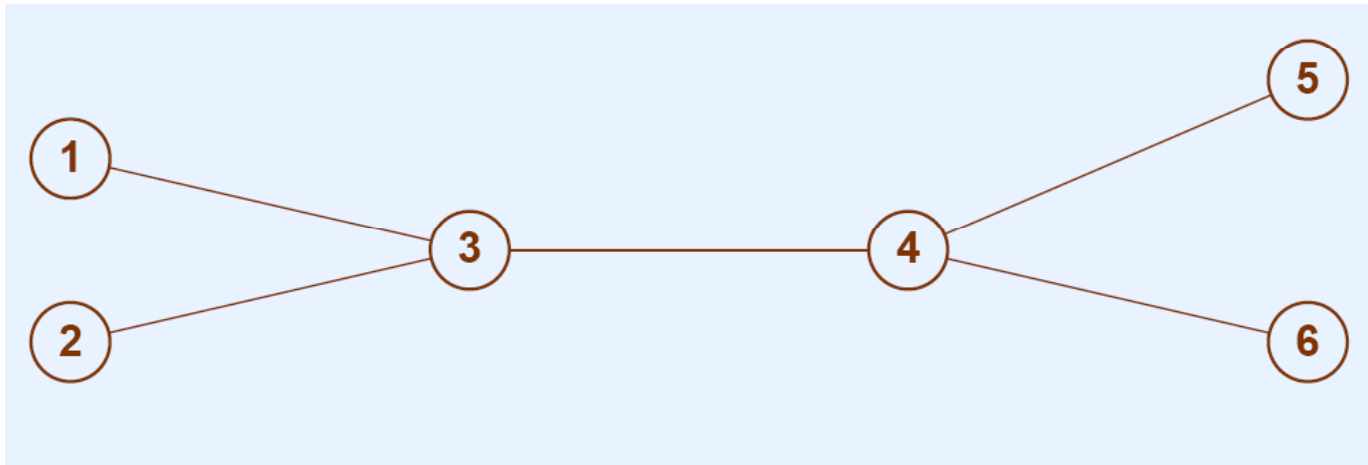
(Window size can't exceed the B)

Congestion

- Fundamental problem in networks
- Caused by traffic, not hardware failure
- Analogous to congestion on a highway
- Principle cause of delay

→ (Bottleneck)

Illustration Of Architecture That Can Experience Congestion



- Multiple sources
- Bottleneck

Congestion And Loss

Modern network hardware works well; most packet loss results from congestion, not from hardware failure.

If entire network becomes unusable → congestion collapse → packet switch (on the way) sets a bit in header → destination acknowledge the sender informing congestion occurred.

Avoiding Congestion

- Rate control
 - Limit rate of data transmission
 - Performed by sending computer
(reducing sliding window size etc.) → balancing needed
 - Performed by network
- Network rate control
 - Monitor incoming traffic
 - Drop or reject packets over rate
 - Called traffic shaping

Summary

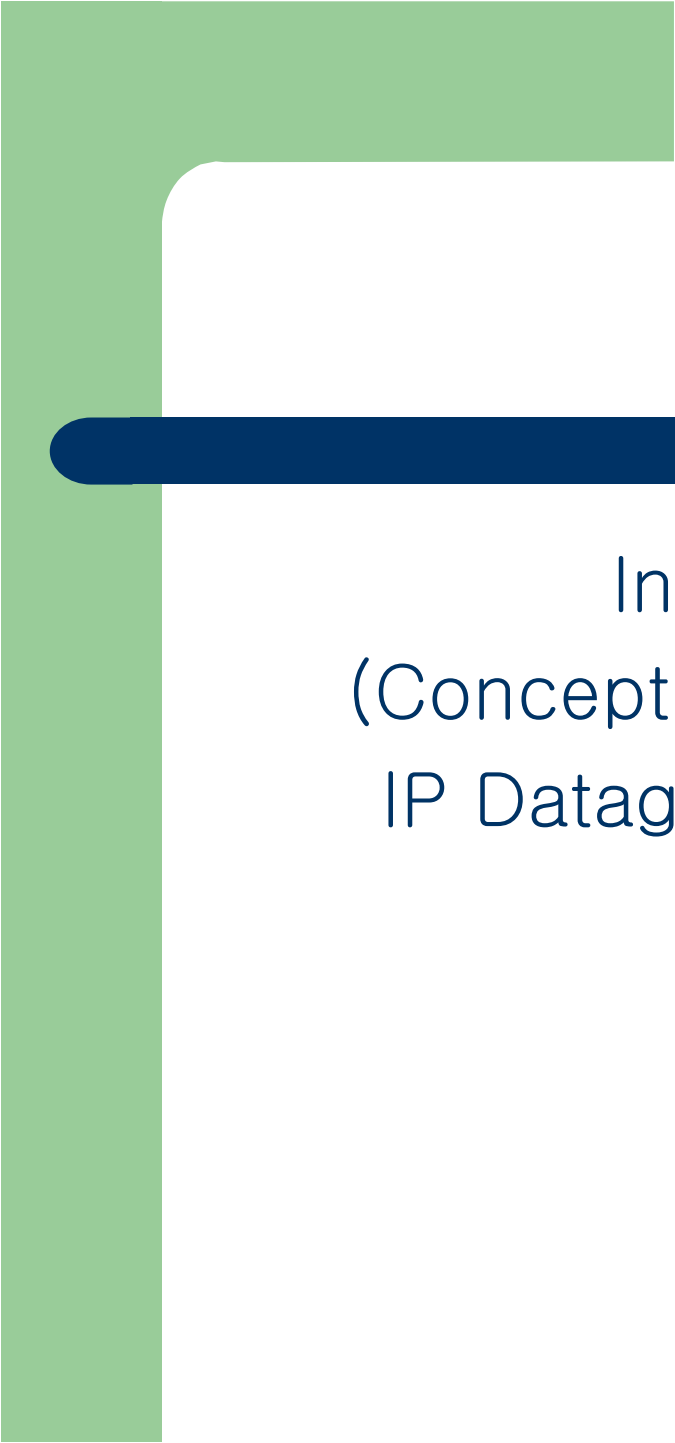
- Protocols
 - Rules for communication
 - Specify syntax (forms) and semantics (meanings)
 - Complex
- Protocol layering
 - Intended for protocol designers
 - Helps organize set of protocols
 - Each layer handles one problem

Summary (continued)

- Problems and techniques
 - Corruption: parity, checksums, CRCs
 - Duplication, out-of-order delivery: sequence numbers
 - Loss: acknowledgement and retransmission
 - Replay: unique ID
 - Congestion: rate control
 - Data overrun: flow control

Summary (continued)

- Two types of flow control
 - Stop-and-go
 - Sliding window
- Sliding window
 - Receiver advertises buffer
 - Sender can fill entire buffer
 - Produces higher performance



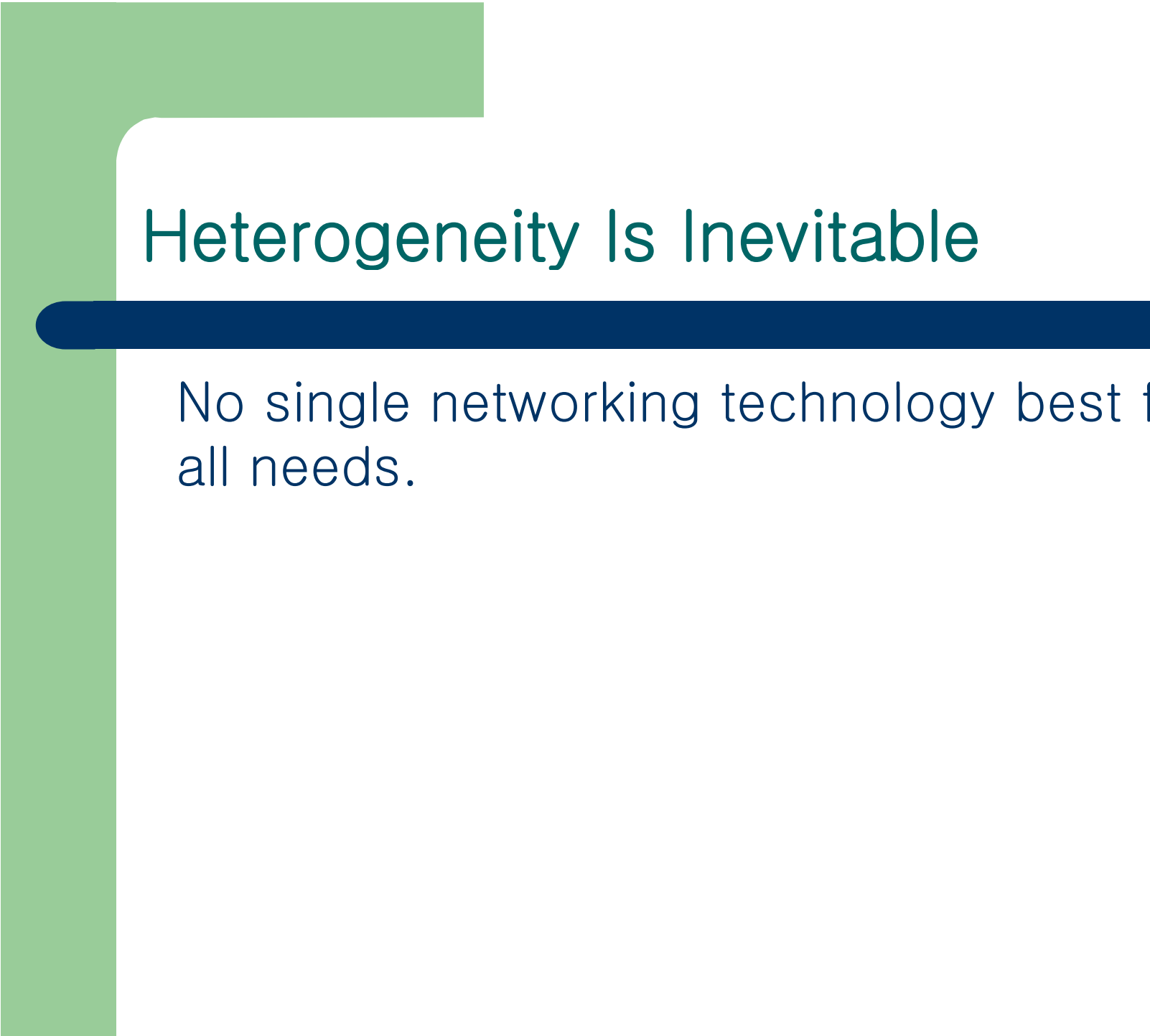
PART X

Internetworking Part 1
(Concept, IP Addressing, IP Routing,
IP Datagrams, Address Resolution)

Motivation For Internetworking

- LANs
 - Low cost
 - Limited distance
- WANs
 - High cost
 - Unlimited distance

Heterogeneity Is Inevitable



No single networking technology best for all needs.

Universal Service

- Fundamental concept in networking
- Pioneered by telephone system
- Arbitrary pair of computers can communicate
- Desirable
- Difficult in a heterogeneous world

Heterogeneity And Universal Service

- Incompatibilities among networks
 - Electrical properties
 - Signaling and data encoding
 - Packet formats
 - Addresses

The Bottom Line

Although universal service is highly desirable, incompatibilities among network hardware and physical addressing prevent an organization from building a bridged network that includes arbitrary technologies.

An Internetworking (= Internet)

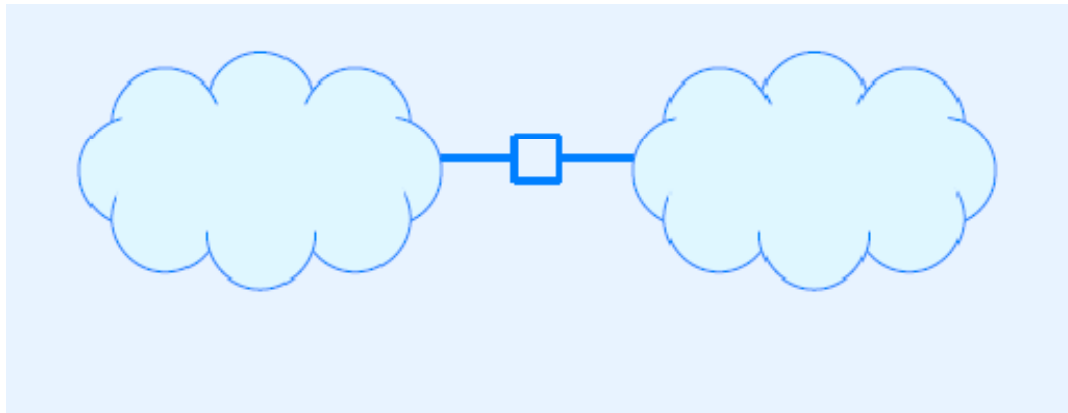
- Begin with heterogeneous network technologies
- Connect the physical networks
- Create software to make resulting system appear homogeneous
- Called an internetwork or internet

Connecting Heterogeneous Networks

- Computer system used
 - Special-purpose
 - Dedicated
 - Works with LAN or WAN technologies
 - Known as
 - * Internet router
 - * Internet gateway

In modern networking parlance, "gateway" and "router" are usually interchangeable and refer to devices with the same purpose. For a network administrator running a large company network, there can be a difference between a router and a gateway. The company may have several subnets connected together via routers, and all of the routers are connected to a gateway device that allows the computers on the subnets to access the Internet. In this case, at a technical level the gateway and the routers perform the same networking functions, but the gateway is an actual "gateway" to the Internet and the routers allow inter-subnet communication – the real difference lies in a device's role in the network structure.

Illustration Of An Internet Router

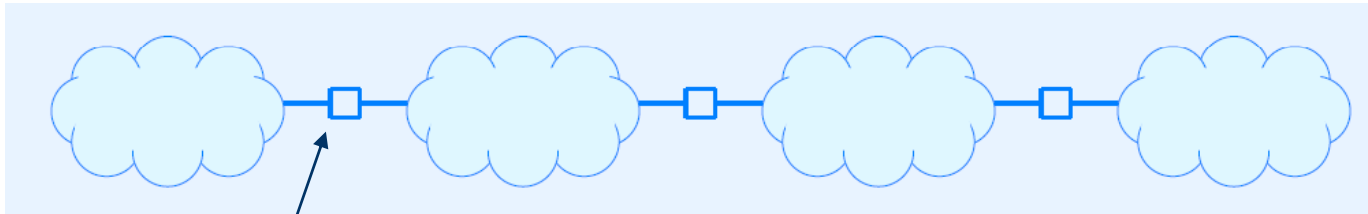


- Cloud denotes arbitrary network technology
- One interface per network
(Router is treated as a computer)

Important Idea

A **router** can interconnect networks that use different technologies, including different media and media access techniques, physical addressing schemes, or frame formats.

Internet Architecture

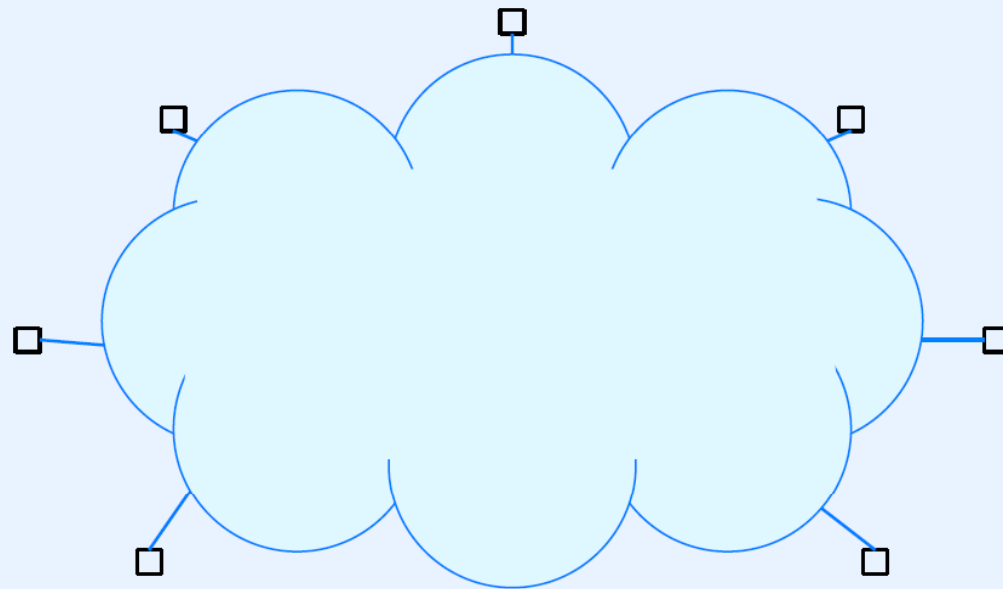


- Multiple
 - Networks
 - Routers interconnecting networks
- Host computer connects to a network
- Single router has insufficient
(to connect all the networks)
 - CPU power and memory
 - I/O capability

Internetworking

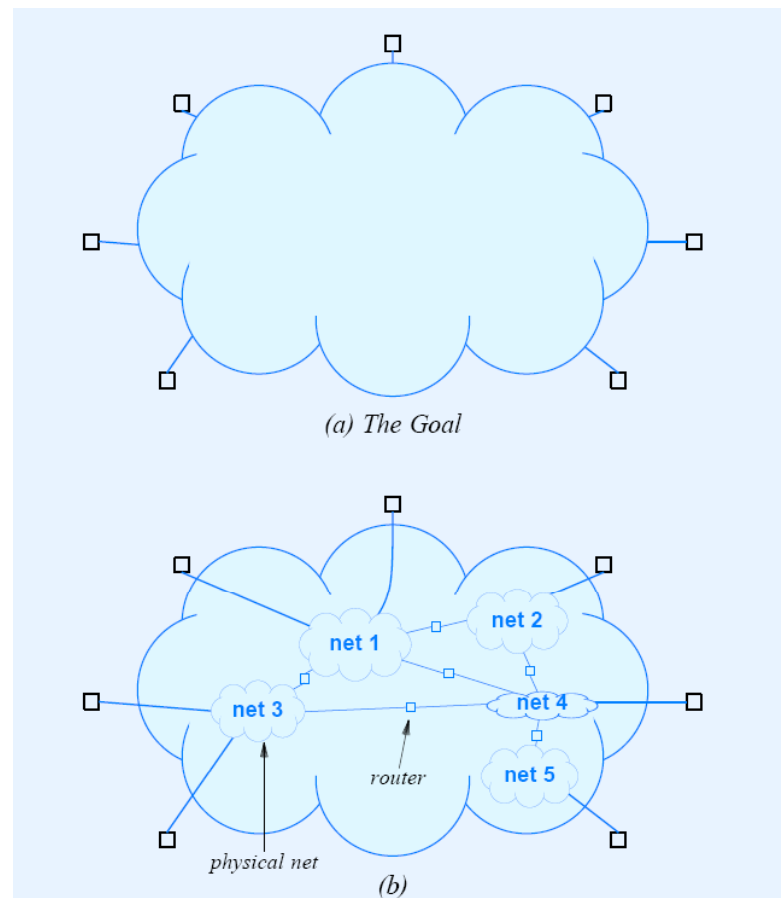
- Goal: communication system
 - Seamless
 - Uniform
 - General-purpose
 - Universal
 - Hides heterogeneity from user

The Internet Concept



(a) The Goal

The Internet Concept



Virtual Network

Heterogeneous network
connected through routers
in reality

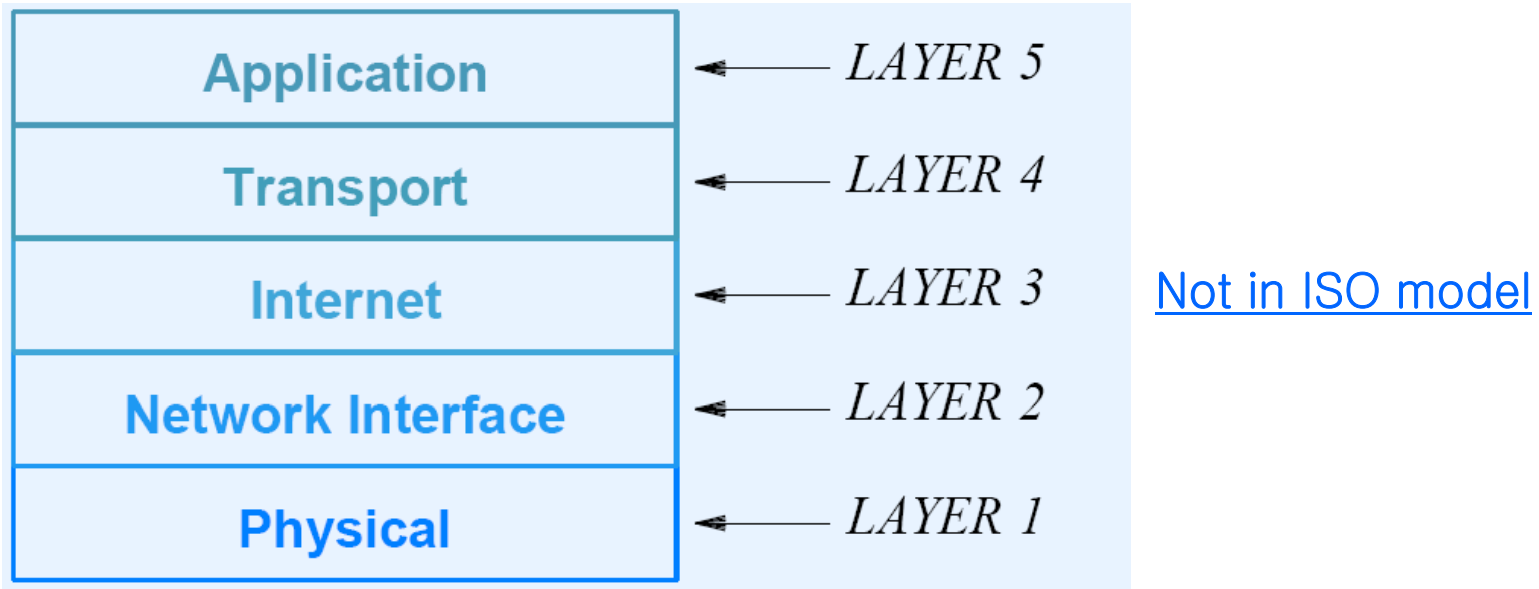
To Hide Heterogeneity

- Create “virtual” network
- Invent
 - Addressing scheme
 - Naming scheme
- Implement with
 - Protocol software
- Note: protocol software needed on both hosts and routers

Internet Protocols

- Known as TCP/IP
- Many protocols comprise suite
- Designed to work together
- Divided into five conceptual layers

Layering Used With TCP/IP



Note: TCP/IP layering model replaces the old ISO model

TCP/IP Layers

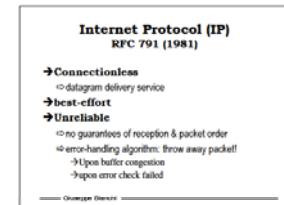
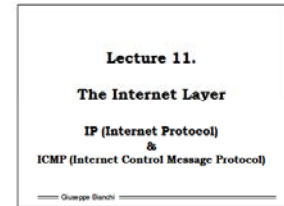
- Layer 1: Physical
 - Basic network hardware
- Layer 2: Network Interface (→ like Data Link layer)
 - MAC frame format
 - MAC addressing
 - Interface between computer and network (NIC)
- Layer 3: Internet
 - Facilities to send packets across internet composed of multiple routers
 - Specify the format of packets
 - How to forward packets through routers

TCP/IP Layers (continued)

- Layer 4: Transport
 - Transport from an application on one computer to application on another → (Reliability)
- Layer 5: Application
 - Everything else

Internet Protocol (IP)

- Only protocol at Layer 3
- Fundamental in suite
- Defines
 - Internet addressing
 - Internet packet format
 - Internet routing



IP Addressing

- Abstraction
- Independent of hardware addressing
(Since Multiple network technologies use different address formats)
- Used by
 - Higher-layer protocols
 - Applications

IP Address

- Virtual
 - Only understood by software
- Used for all communication
- 32-bit integer
- Unique value for each host